RANDOMIZATION in SYSTEMS and CONTROL: a CHANGE of PERSPECTIVE

Marco C. Campi

with ...

#### Simone Garatti

#### Giuseppe Calafiore

Algo Care'



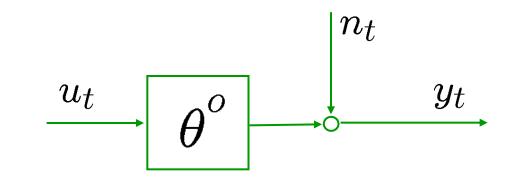


Maria Prandini

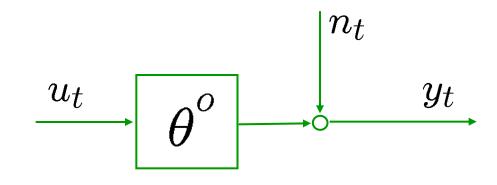
Erik Weyer

#### outline:

- (i) focusing on randomization: an example in estimation theory
- (ii) what is a randomized algorithm?
- (iii) where are we?
- (iv) uncertain systems



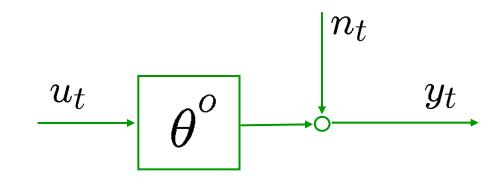
$$y_t = \theta^o u_t + n_t$$



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• select  $u_t, t = 1, \ldots, N$ 

• measure  $y_t$ 



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**<u>goal</u>**: provide an interval  $\Theta$  for  $\theta^o$ M:  $(u_t, y_t) \rightarrow \Theta$ 

requirements:

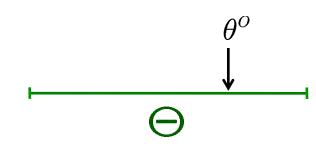
requirements:

(i) no assumptions on  $n_t$ 

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(ii) we want to issue a certificate of reliability on  $\Theta$  valid  $\forall n_t$ 



is this at all possible?

- take an input sequence  $u_t$
- measure  $y_t$

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$$\begin{array}{c} \theta^{o} = 3 \\ u_{t} \\ y_{t} \end{array} \end{array} \xrightarrow{M} \begin{array}{c} 1.9 \\ \bigcirc \end{array} \begin{array}{c} 2.5 \\ \bigcirc \end{array} \end{array}$$

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$$\begin{array}{c} u_t \\ y_t \end{array} \right\} \xrightarrow{M} 1.9 \underbrace{2.5}_{\bigcirc} \psi \\ y_t = 3u_t + n_t \\ \psi \\ n_t = -3u_t + y_t \end{array}$$

- a deterministic algorithm that comes with a certificate of reliability does not exist!  a deterministic algorithm that comes with a certificate of reliability does not exist!

- shall we give up?

moving one step ahead ... a change of perspective moving one step ahead ... a change of perspective

**Algorithm** 

1) 
$$u_t = \begin{cases} +1, & \text{with prob} = 0.5 \\ -1, & \text{with prob} = 0.5 \end{cases}$$
 (R. Fisher)

moving one step ahead ... a change of perspective

#### **Algorithm**

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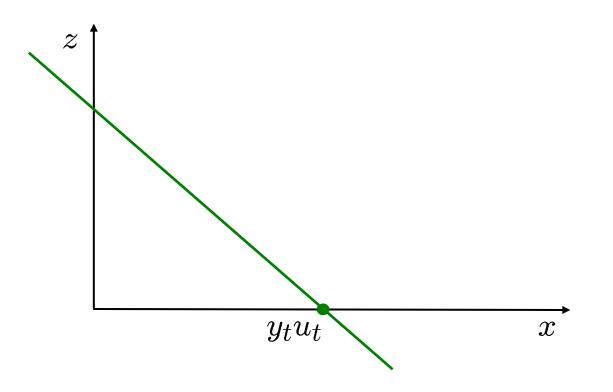
2) 
$$M: (u_t, y_t) \to \Theta$$
 needs some explanation

(work done with E. Weyer)





 $z = y_t \cdot u_t - x$  -45 degrees line  $t = 1, 2, \dots, 7$ 





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construct some averages:

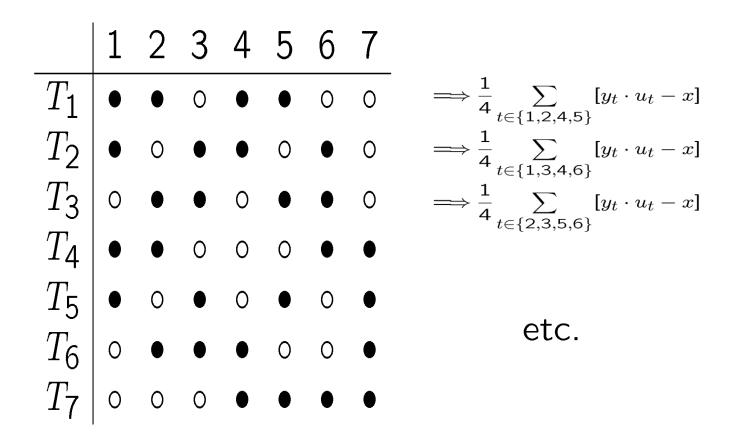
$$\frac{1}{|T_1|} \sum_{t \in T_1} [y_t \cdot u_t - x]$$
$$\frac{1}{|T_2|} \sum_{t \in T_2} [y_t \cdot u_t - x]$$



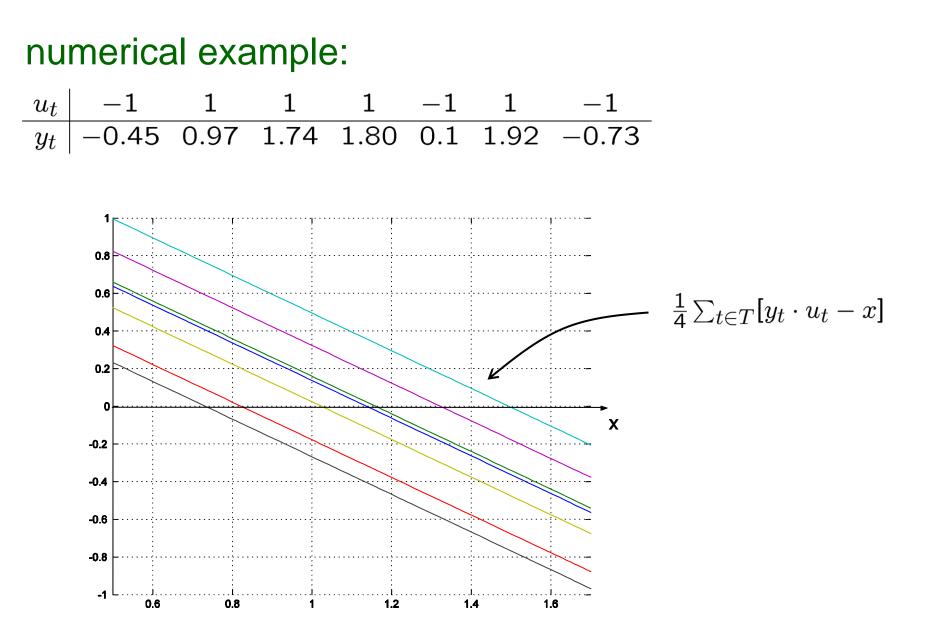
numerical example:



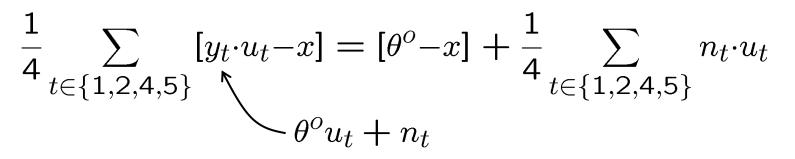
umerical example: $u_t$ -111-11-1 $y_t$ -0.450.971.741.800.11.92-0.73

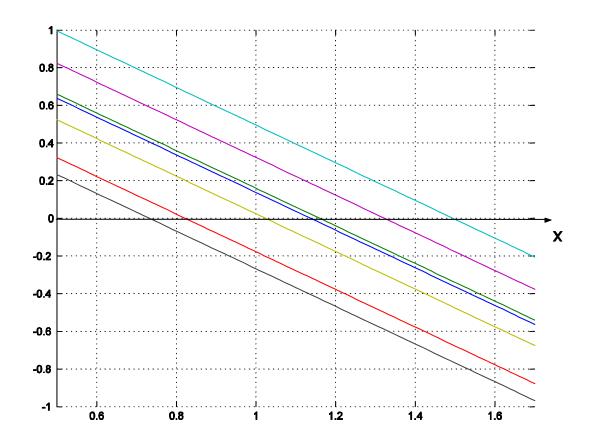


#### $\underline{explaining} M$

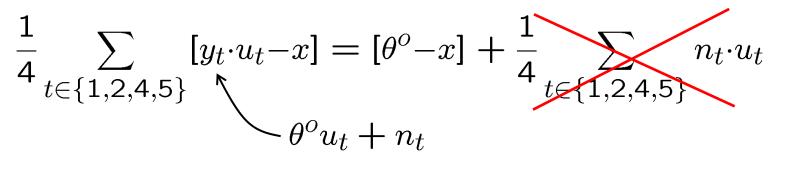


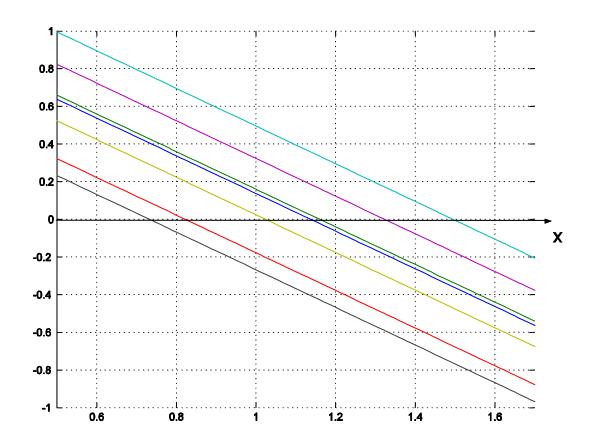
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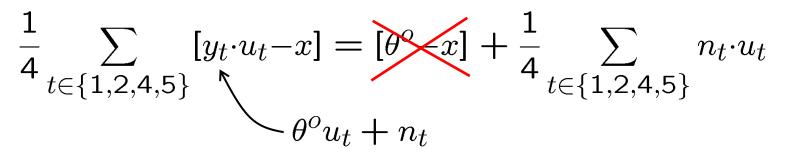


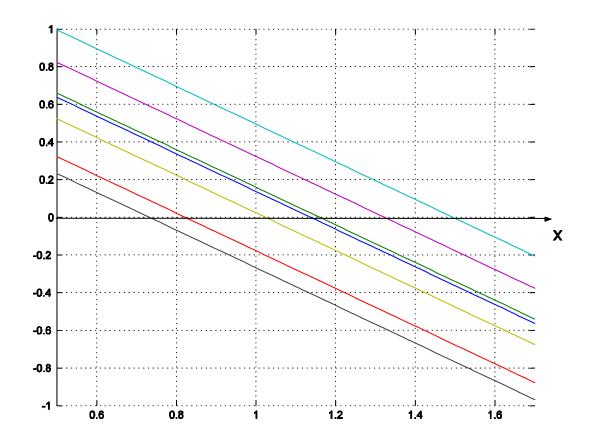
## $\underline{explaining} M$





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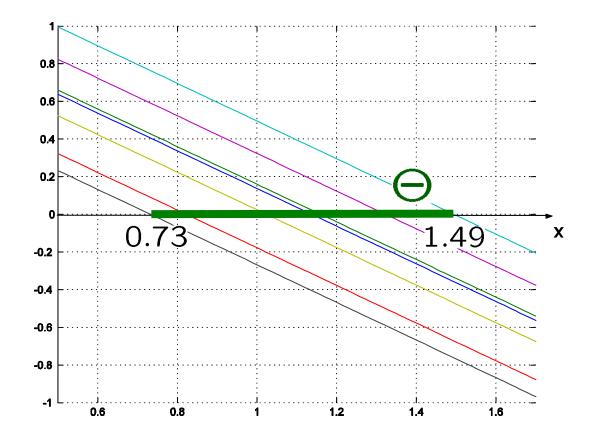




## $\underline{\text{explaining}} \ M$

$$\frac{1}{4} \sum_{t \in \{1,2,4,5\}} [y_t \cdot u_t - x] = [\theta^o - x] + \frac{1}{4} \sum_{t \in \{1,2,4,5\}} n_t \cdot u_t$$

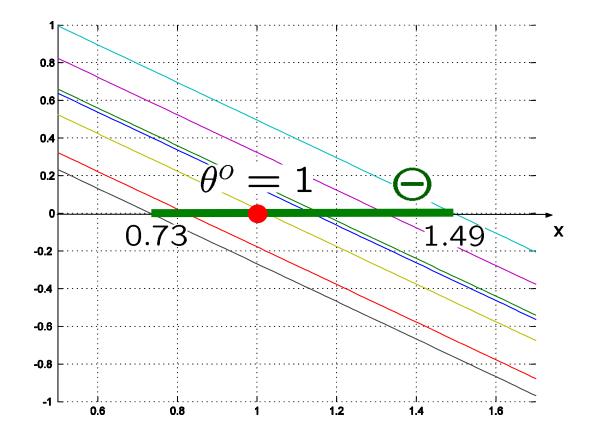
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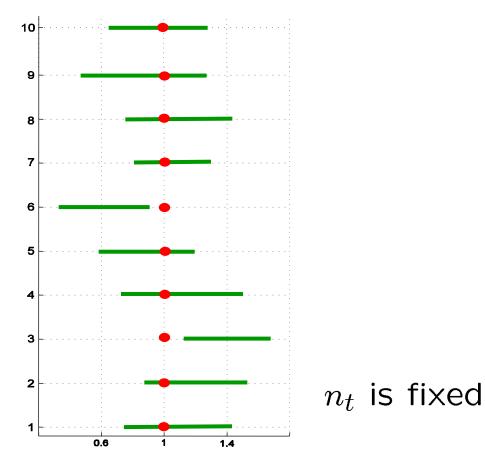
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N = 1023 $\Theta = [0.945, 1.033]$  with probability = 0.975 Theorem (with E. Weyer)

# $\forall n_t, \ \theta^o \in \Theta$ with probability 0.75.

Theorem (consistency)

Under general assumptions,  $\Theta \rightarrow \{\theta^o\}$  as  $N \rightarrow \infty$ 

## in summary:

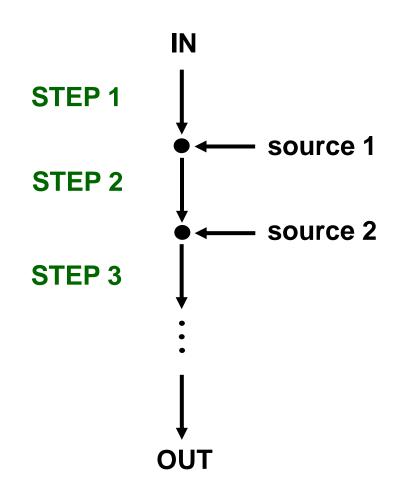
for any  $n_t$ :

- (i)  $\theta^o \in \Theta$  with a precise probability
- (ii) the size of  $\Theta$  depends on the strength of the noise
- (iii)  $\Theta$  shrinks around  $\theta^o$  as *N* increases

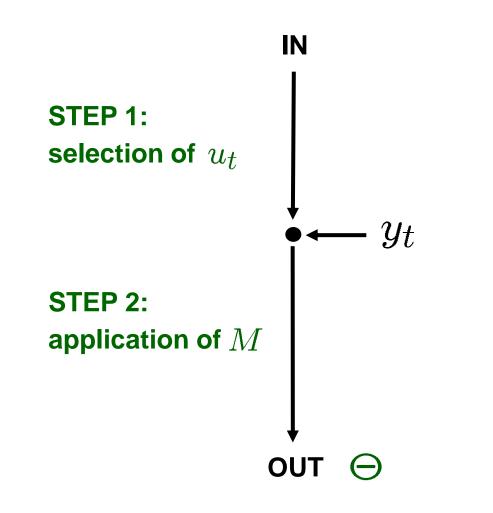
how is all this possible?

# what is a randomized algorithm?

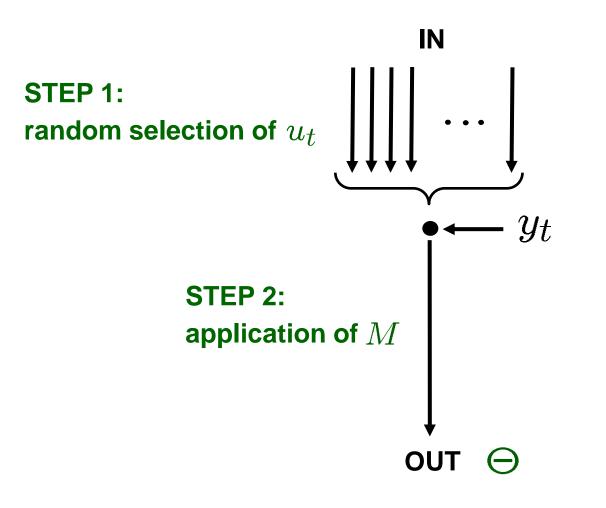
# deterministic algorithm



# deterministic algorithm



# randomized algorithm



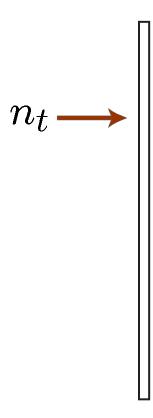
a randomized algorithm is an algorithm where one or more steps are based on a random choice

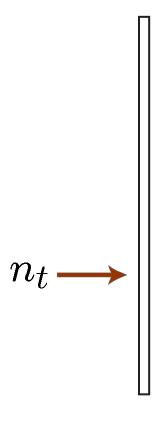
that is – among many deterministic choices – one choice is selected at random according to a probability *P* 

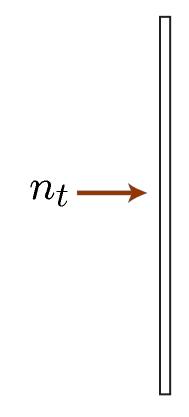
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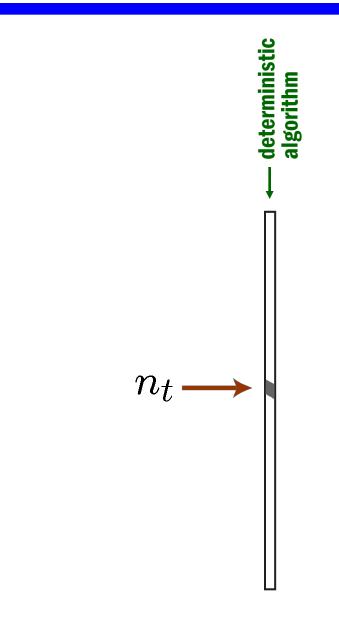
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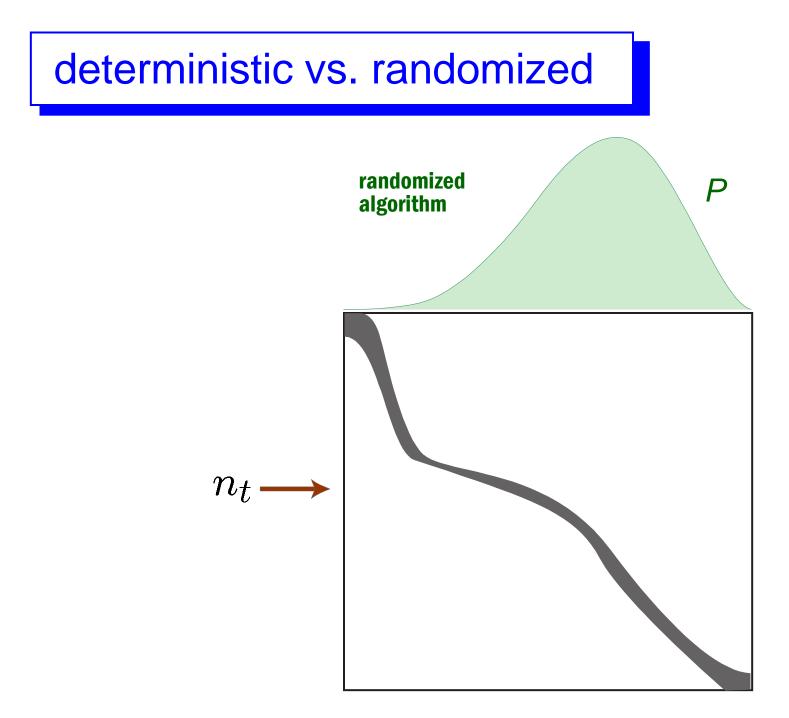
why is this useful?

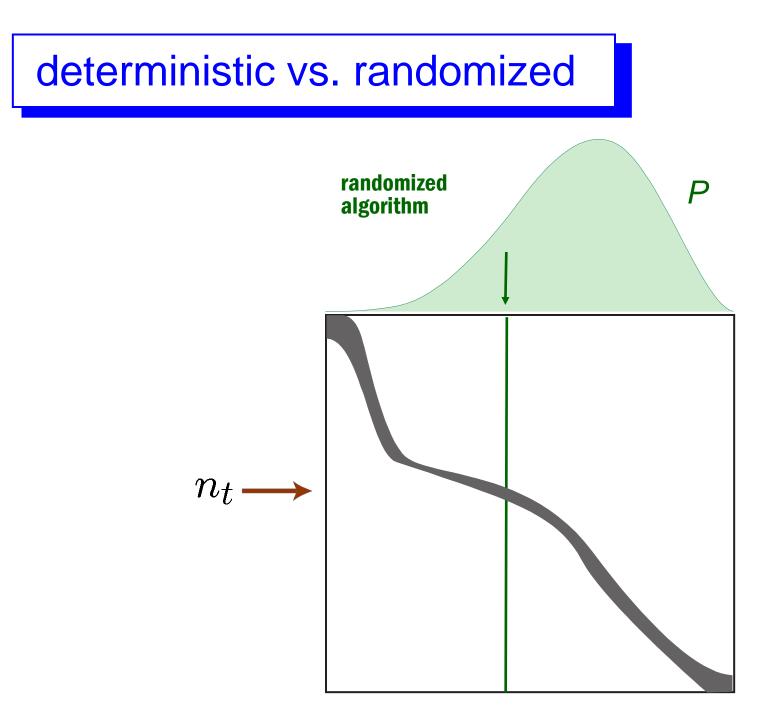


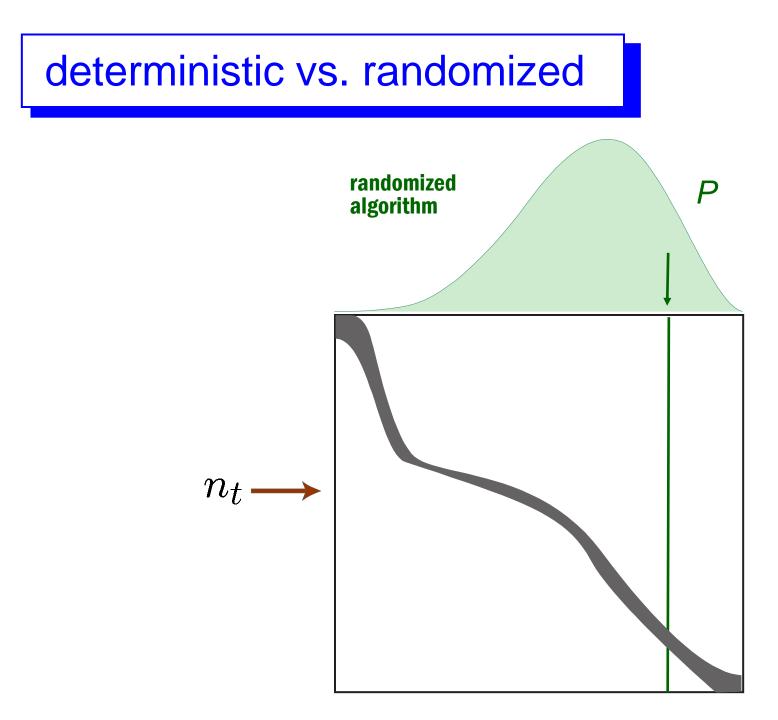


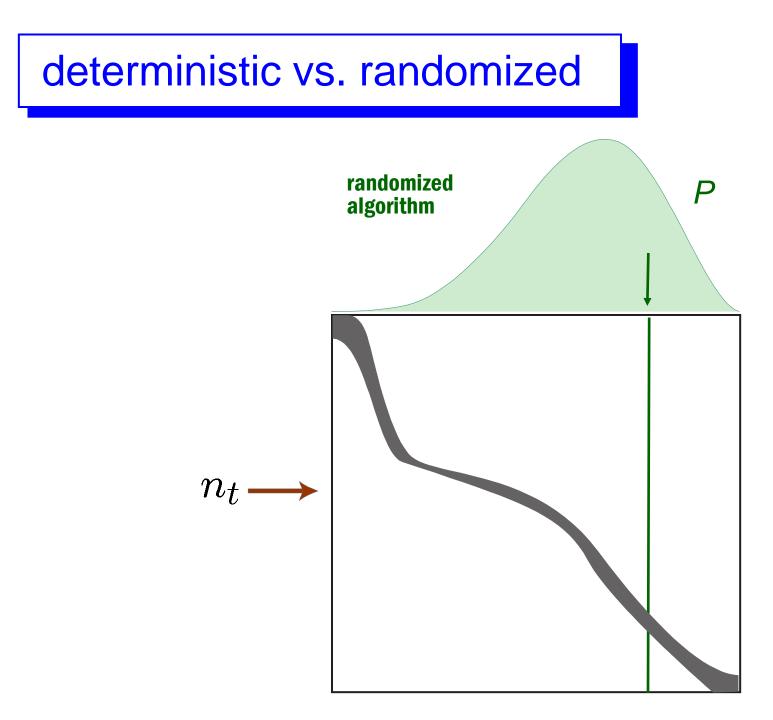


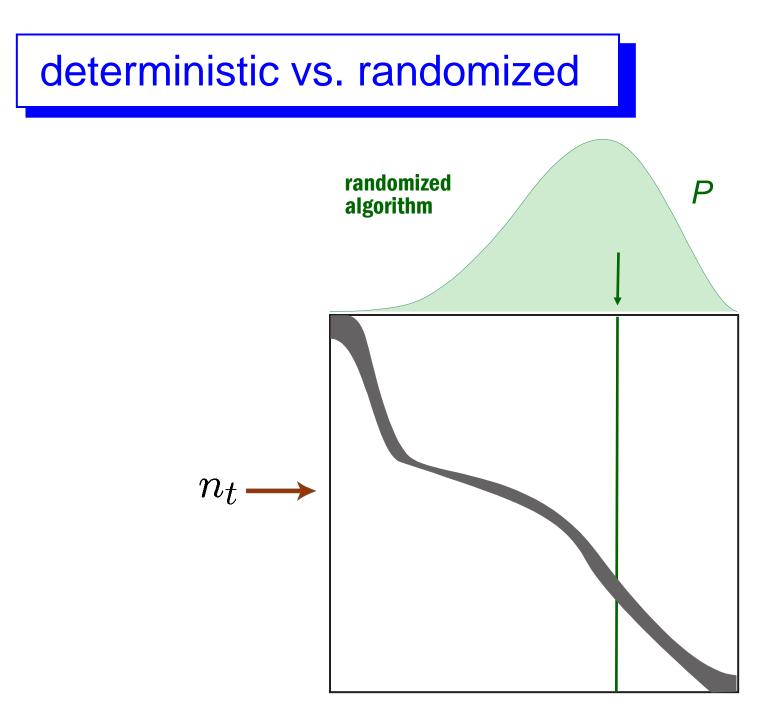


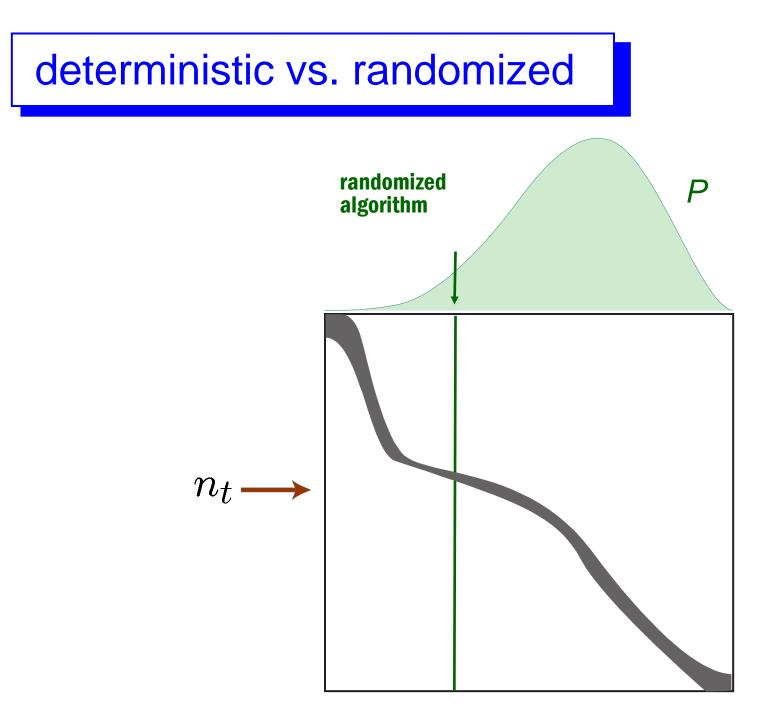


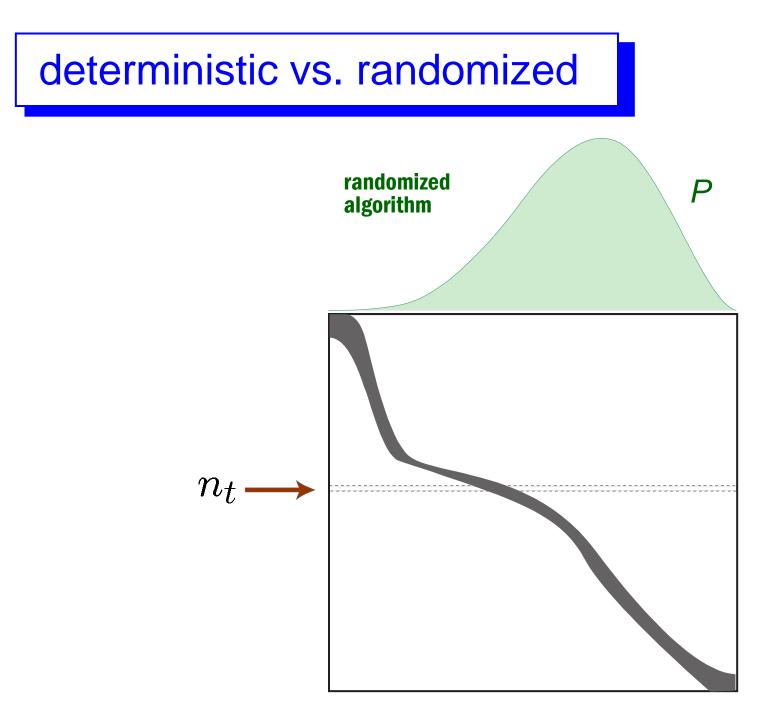












# successful algorithm

successful algorithm

Definition 1 (successful deterministic algorithm) An algorithm is successful if, in all situations, it provides a correct answer.

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... a change of perspective:

Definition 2 (probabilistically successful algorithm) An algorithm is successful with probability *p* if, in all situations, its probability to provide a correct answer is at least *p*. Definition 1 (successful deterministic algorithm) An algorithm is successful if, in all situations, it provides a correct answer.

... a change of perspective:

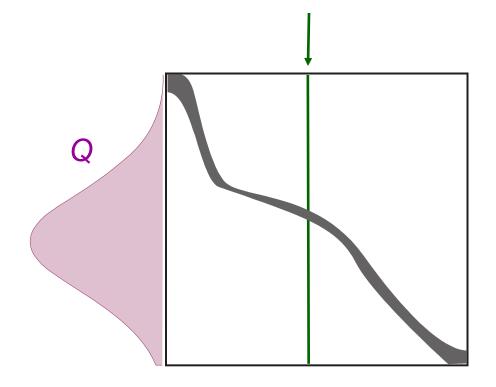
Definition 2 (probabilistically successful algorithm) An algorithm is successful with probability *p* if, in all situations, its probability to provide a correct answer is at least *p*.

this offers an extraordinary opportunity to satisfactorily solve "hopeless" problems



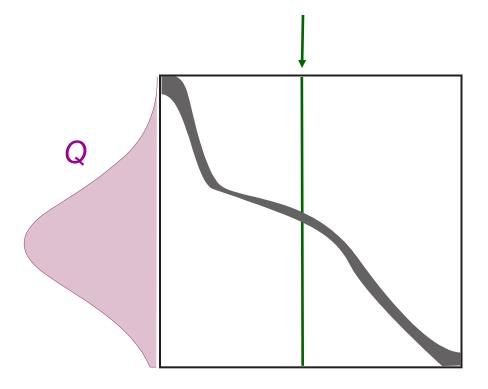
• P exists in the algorithm

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**Bayesian perspective** 

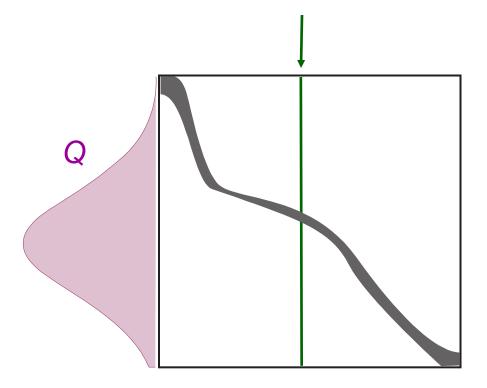
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**Bayesian perspective** 

(1) the result holds with high probability with respect to the situation of application

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**Bayesian perspective** 

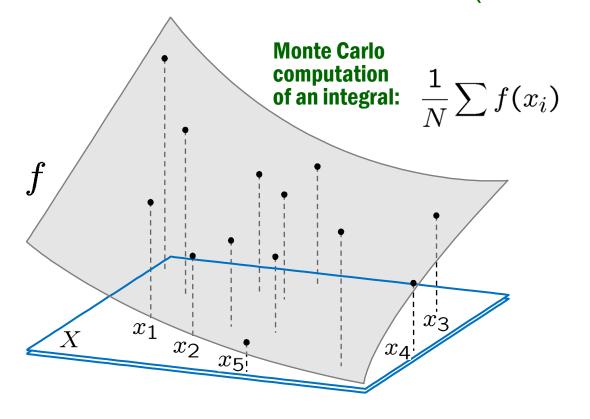
- (1) the result holds with high probability with respect to the situation of application
- (2) Q describes reality  $\rightarrow$  poor modeling



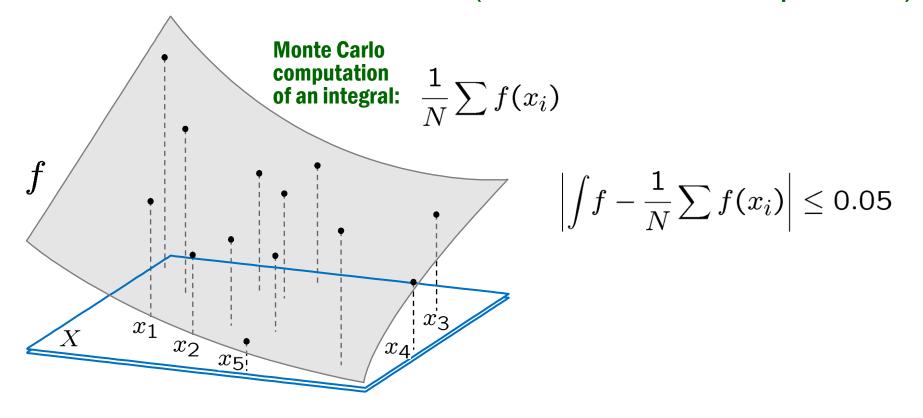
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- often the chance of failure can be made very small (concentration inequalities)

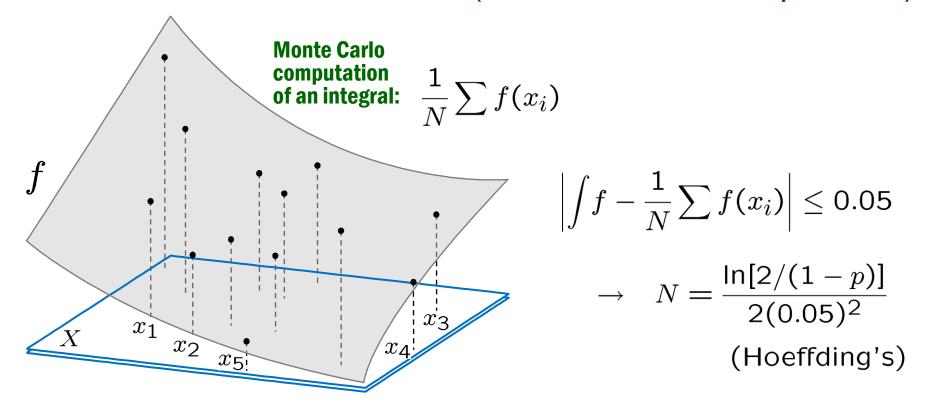
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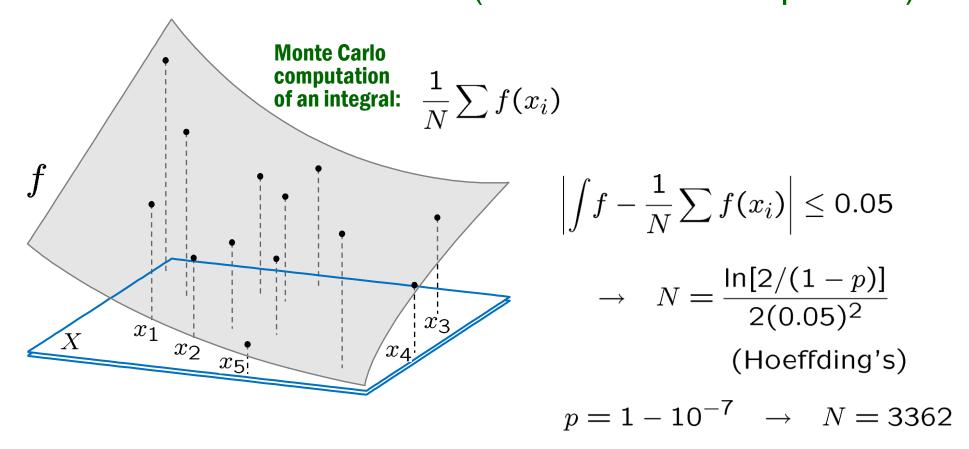
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- often the chance of failure can be made very small (concentration inequalities)
- amazingly powerful results from probability theory can be used to assess the probability of success



## where are we?

- Monte Carlo method (1949)
- computer science
  - sorting
  - counting
  - incremental geometric constructions
  - etc.
- optimization
  - simulating annealing
  - genetic methods
  - large-scale convex optimization
  - etc.
  - Iittle in systems and control

### where are we in systems and control?

#### uncertain systems

- aerospace
- adaptive control
- network control

-etc.

#### • many contributors:

T. Alamo, E.W. Bai, B.R. Barmish, T. Basar, G. Calafiore, A. Chaouki, F. Dabbene, B. De Shutter, L. El Ghaoui, M. Fu, Y. Fujisaki, S. Garatti, H. Ishii, S. Kanev, H. Kimura, C. Lagoa, S.P. Meyn, Y. Oishi, B. Polyak, M. Prandini, P. Shcherbakov, J. Spall, R.F. Stengel, M. Sznaier, V.B. Tadic, R. Tempo, B. Van Roy, M. Verhagen, M. Vidyasagar, K. Zhou

 R. Tempo, G. Calafiore, and F. Dabbene (2005). "Randomized algorithms for analysis and control of uncertain systems". Springer-Verlag. uncertain systems: design

### robust min-max design (convex)

uncertain systems: design

#### robust min-max design (convex)

(non-convex — Alamo, Camacho, Tempo)

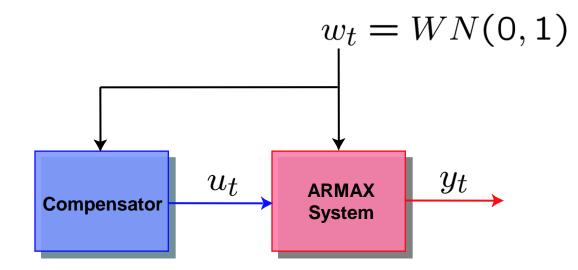
a successful story of randomization: robust min-max convex design a successful story of randomization: robust min-max convex design

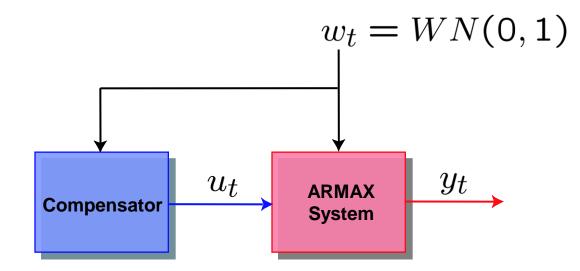
$$\min_{\theta} \left[ \max_{\delta \in \Delta} \ell(\theta, \delta) \right]$$

a successful story of randomization: robust min-max convex design

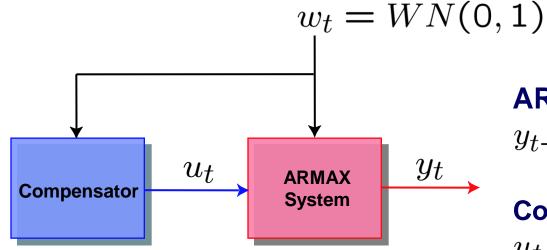
$$\min_{\theta} \left[ \max_{\delta \in \Delta} \ell(\theta, \delta) \right]$$

#### robust min-max design is hard!





Objective: reduce the effect of noise on y

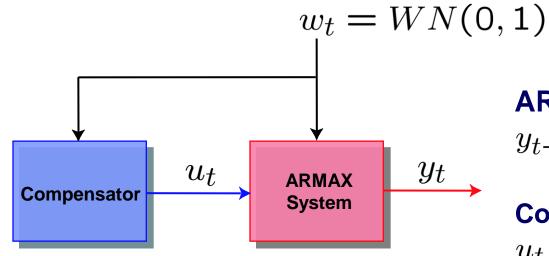


#### **ARMAX System:**

 $y_{t+1} = ay_t + bu_t + cw_t + dw_{t-1}$ 

Compensator:  $u_t = k_1 w_t + k_2 w_{t-1}$ 

Goal: min  $var[y_t]$ 



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$$var[y_t] = \frac{(c+bk_1)^2 + (d+bk_2)^2 + 2a(c+bk_1)(d+bk_2)}{1-a^2}$$

system parameters unknown:  $a, b, c, d \in \Delta$ 

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#### even a problem as simple as this is difficult for a generic $\Delta$

# other problems in robust control

- state-feedback stabilization
- $H_{\infty}$  control
- *H*<sub>2</sub> control
- LPV control
  - :
  - •

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- state-feedback stabilization
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- H<sub>2</sub> control
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  - :
  - •
- ... and systems theory
- model reduction
- prediction
  - :

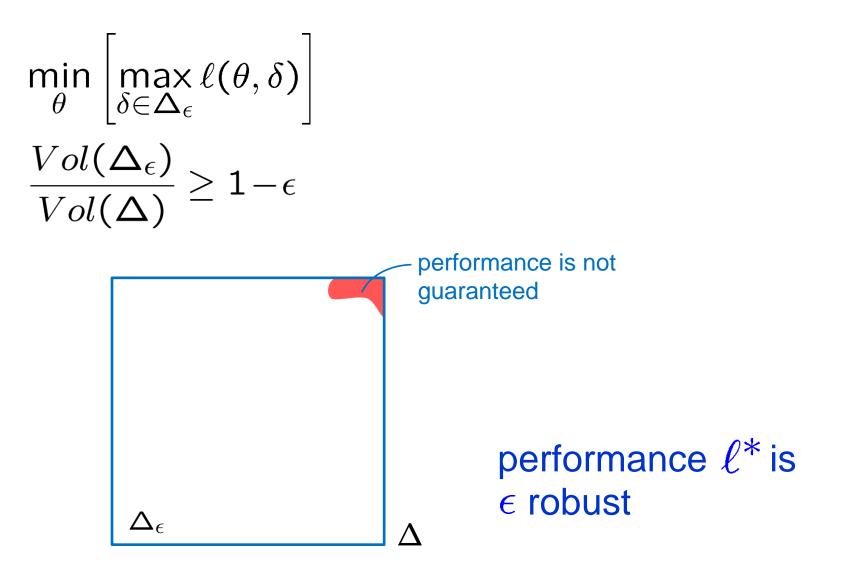
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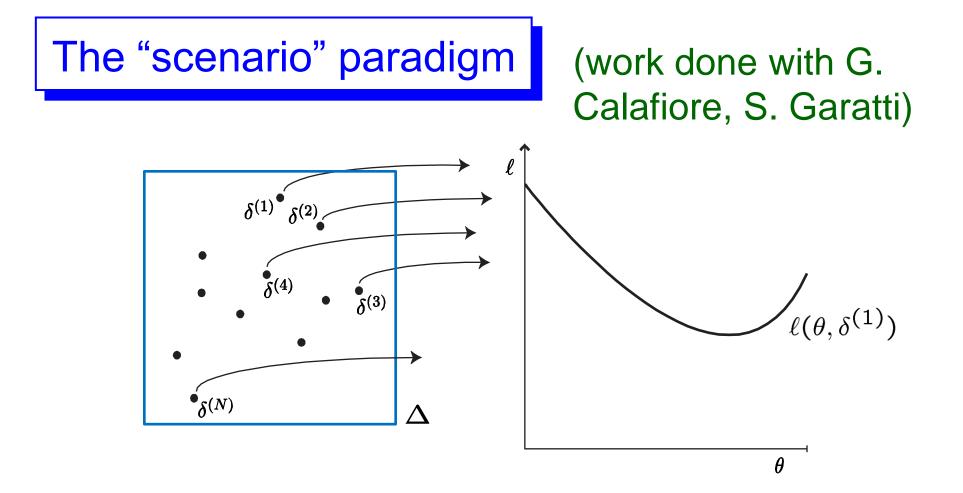
- robust min-max design is hard!
- what can we do?

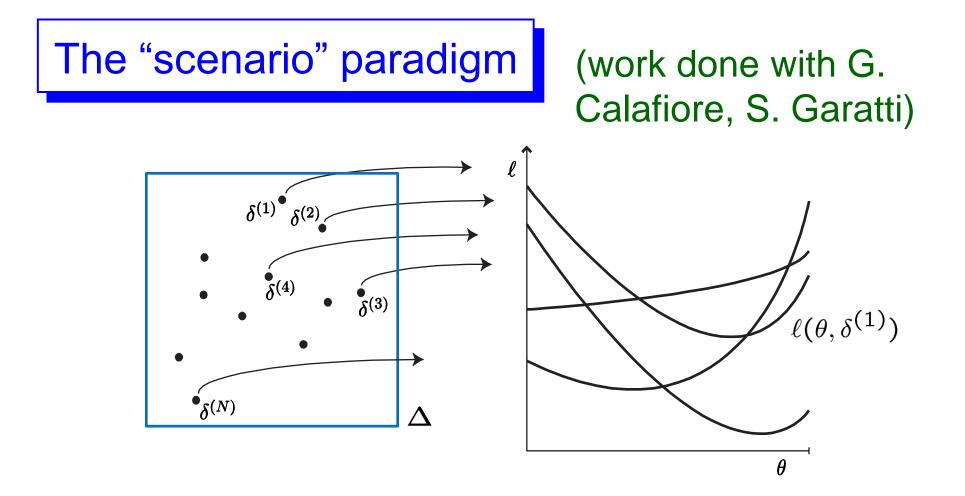
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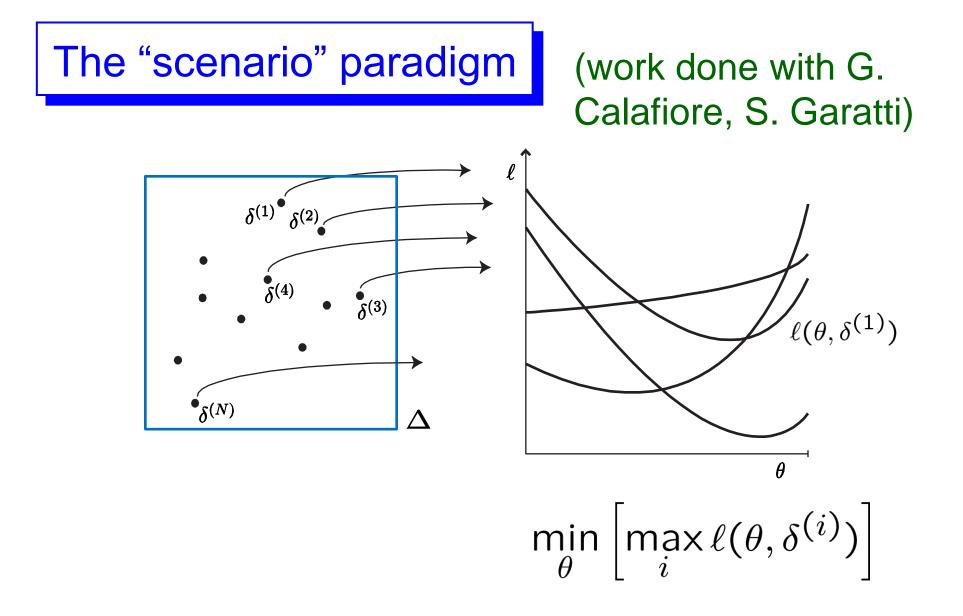
- we need to accept a compromise

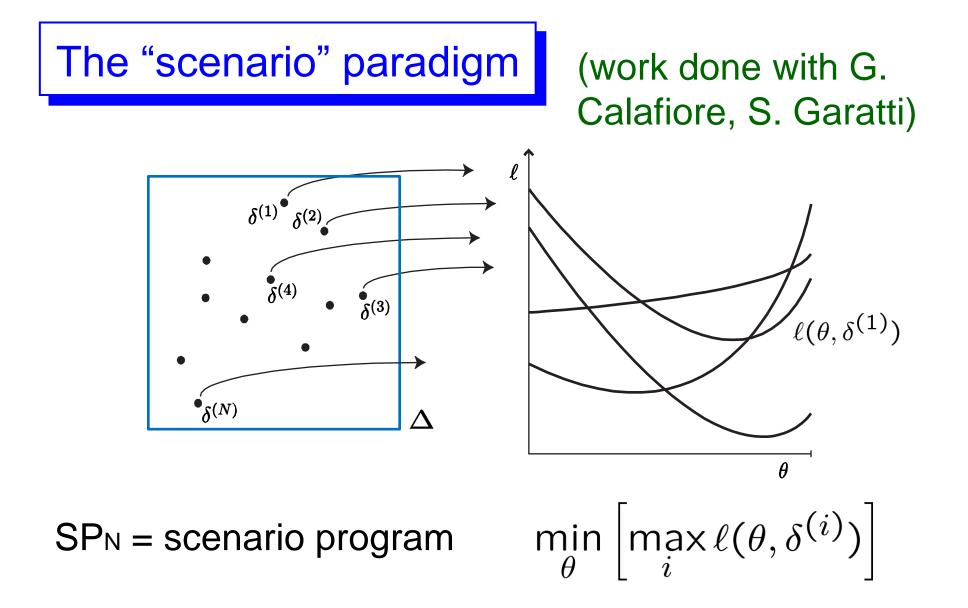
chance-constrained optimization

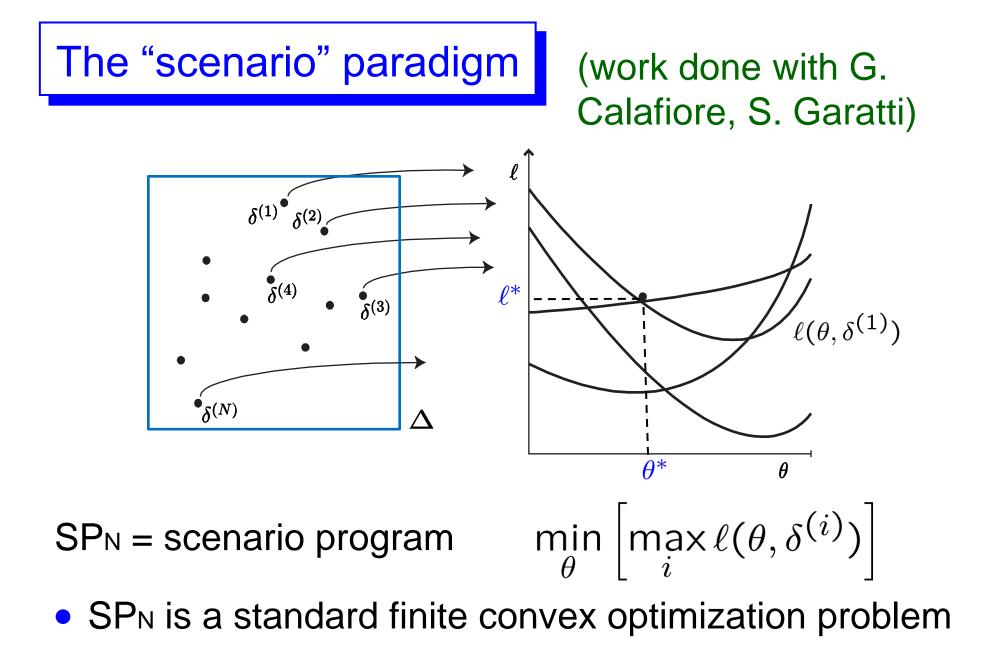


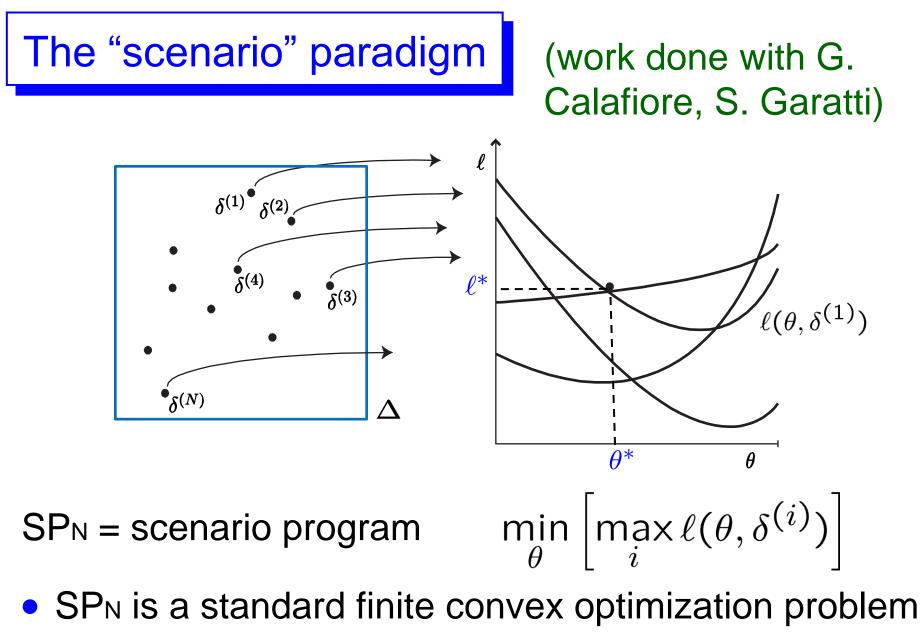












•  $\ell^*$  is superoptimal

# Fundamental question:

How robust is  $\ell^*$ ?

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from the "visible" to the "invisible"

### Theorem (with S. Garatti – G. Calafiore)

Fix  $p \in (0, 1)$  (probability of success)

If 
$$N = \frac{2}{\epsilon} \left( \ln \frac{1}{1-p} + n_{\theta} \right)$$
, then,

with probability of success p,  $\ell^*$  is  $\epsilon$  robust.

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$$p = 1 - 10^{-7} \quad \rightarrow \quad \ln \frac{1}{1-p} \sim 16$$

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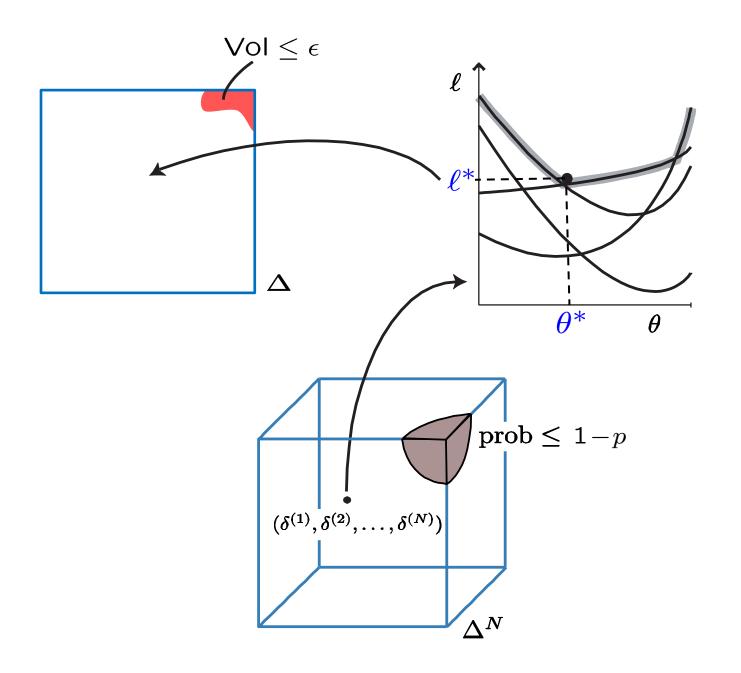
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applicable to all convex problems!



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$$\Delta = \{a, b, c, d : a = 0.45 + 0.5 * (1 - e^{-8 \cdot 10^3 (\sigma_1^2 + \sigma_2^2)}), \\b = 1 + \sigma_2^2, \\c = 0.2 + (\sigma_2 + \sin(\sigma_2) + 0.1) \cdot \sin(2\pi\sigma_2), \\d = 0.5 + \sigma_1^2 \cos(\sigma_2),$$

$$(\sigma_1, \sigma_2) \in [-1/3, 1/3]^2 \}.$$

$$\varepsilon = 0.01 \quad p = 1 - 10^{-7} \quad \Longrightarrow \quad N = 2158$$

sample:  $a_i, b_i, c_i, d_i \in \Delta$ , i = 1, 2, ..., 2158;

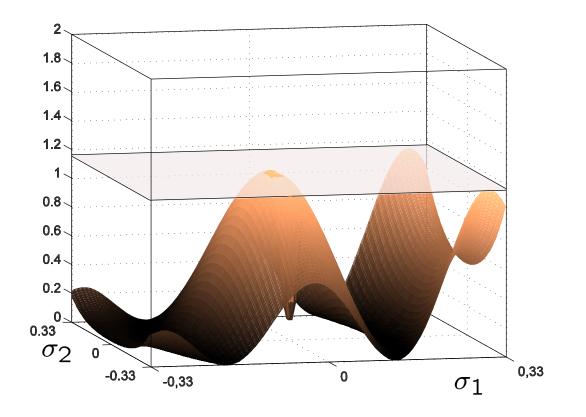
#### solve:

$$\min_{k_1,k_2} \left[ \max_i \frac{(c_i + b_i k_1)^2 + (d_i + b_i k_2)^2 + 2a_i (c_i + b_i k_1) (d_i + b_i k_2)}{1 - a_i^2} \right]$$

$$\implies k_1^* = -0.50, \quad k_2^* = -0.53, \quad \ell^* = 1.16$$

# $\ell^* = 1.16$ $\longrightarrow$ Output variance below 1.16 for all plants but a small fraction (1%)

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once more, a tough problem has turned into a solvable one through randomization, ... provided we accept an  $\epsilon$  risk



- randomization changes our <u>perspective of problem</u> <u>solvability</u>

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- the probability of success depends on an <u>artificial</u> *P* and can be assessed with <u>extraordinarily powerful</u> <u>probabilistic tools</u>

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- can prove <u>useful in many more problems</u> in systems and control, especially at the boundary of control, communication and computation THANK YOU

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